Introduction to Application Performance Analysis with CrayPat

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Cray Inc.

Performance Optimization

We want to get the most science through a supercomputing system as possible

The more efficient codes are the more productive scientists and engineers can be 90

101

Performance Optimization

- Adapting the problem to the underlying hardware
- Combination of many aspects
 - Effective algorithms
 - Implementation: Processor utilization & efficient memory use
 - Parallel scalability
- Important to understand interactions
 - Algorithm code compiler libraries hardware
- Performance is not portable!

Performance analysis

To optimise code we must know what is taking the time



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Not going to touch the source code?

- Find the *compiler* and its *compiler flags* that yield the best performance
- Employ *tuned libraries* wherever possible
- Find suitable settings for *environment parameters*
- Mind the I/O
 - Do not checkpoint too often
 - Do not ask for the output you do not need

Why does scaling end?

- Amount of data per process small computation takes little time compared to communication
- Amdahl's law in general
 - E.g., single-writer or stderr I/O
- Load imbalance
- Communication that scales badly with N_{proc}



Application timing

Most basic information: total wall clock time

- Built-in timers in the program (e.g. MPI_Wtime)
- System commands (e.g. time) or batch system statistics

Built-in timers can provide also more fine-grained information

- Have to be inserted by hand
- Typically, no information about hardware related issues e.g. cache utilization
- Information about load imbalance and communication statistics of parallel program is difficult to obtain

Performance analysis tools

Instrumentation of code

- Adding special measurement code to binary
 - Special commands, compiler/linker wrappers
 - Automatic or manual
- Normally all routines do not need to be measured

• Measurement: running the instrumented binary

- Profile: sum of events over time
- Trace: sequence of events over time

Analysis

- Text based analysis reports
- Visualization

Sampling

Advantages

- Only need to instrument main routine
- Low Overhead depends only on sampling frequency
- Smaller volumes of data
 produced

Disadvantages

- Only statistical averages available
- Limited information from performance counters

Event Tracing

Advantages

- More accurate and more detailed information
- Data collected from every traced function call not statistical averages

Disadvantages

- Increased overheads as number of function calls increases
- Huge volumes of data generated

Guided tracing = trace only program parts that consume a significant portion of the total time In Cray Performance Analysis Toolkit this is referred to as COMPUTE | STORE |"automatic²profiling analysis"(APA)

Step 1: Choose a test problem

- The dataset used in the analysis should
 - Make scientific sense, i.e. resemble the intended use of the code
 - Be large enough for getting a good view on scalability
 - Be runable in a reasonable time
 - For instance, with simulation codes almost a full-blown model but run only for a few time steps
- Should be run long enough that initialization/finalization stages are not exaggerated
 - Alternatively, we can exclude them during the analysis

Step 2: Measure Scalability

- Run the uninstrumented code with different core counts and see where the parallel scaling stops
- Usually we look at strong scaling
 - Also weak scaling is definitely of interest

here?

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Step 3: Run instrumented version of application

- Obtain first a sampling profile to find which user functions should be traced
 - With a large/complex software, one should not trace them all: it causes excessive overhead
- Make an instrumented exe with tracing time-consuming user functions plus e.g. MPI, I/O and library (BLAS, FFT,...) calls
- Execute and record the first analysis with
 - The core count where the scalability is still ok
 - The core count where the scalability has ended
 - and identify the largest differences between these profiles
- CrayPat has an Automatic Profile Analysis (APA) mode to handle this process:

Steps to Collect Performance Data

- Access performance tools software
 - module load perftools-base
 - module load perftools-lite
- Build instrumented version of the application keeping .o files (CCE: -h keepfiles)
 - make clean
 - make
 - You should get an instrumented version program a.out
 - This has been instrumented for sampling (automatic profiling analysis), check with
 - strings a.out | grep 'CrayPat/X' CrayPat/X: Version 6.3.0 Revision 14319 09/02/15 13:51:12

• Run application to get top time consuming routines

- aprun … a.out (or qsub <pat script>)
- You should get *.rpt and a *.ap2 files
- The report in *.rpt is additionally printed to stdout

Example: Sampling report



Example: Sampling report (2)

. . .

. . .

When the use of a shared resource like memory bandwidth is unbalanced across nodes, total execution time may be reduced with a rank order that improves the balance. The metric used here for resource usage is: USER Samp

For each node, the metric values for the ranks on that node are summed. The maximum and average value of those sums are shown below for both the current rank order and a custom rank order that seeks to reduce the maximum value.

A file named MPICH_RANK_ORDER.USER_Samp was generated along with this report and contains usage instructions and the Custom rank order from the following table.

Rank	Node Re	duction	Maximum	Avera
Order	Metric	in Max	Value	Value
	Imb.	Value		

Table 2: I	File Input St	tats by Filena	ame		
Read	Read	Read Rate	Reads	Bytes/	File Name[max15]
Time	MBytes	MBytes/sec		Call	PE=HIDE
0.113291	0.544238	4.803892	2,964.0	192.54	Total
		2 751054	1 596 0	1 4 1 7 0	tenelegy fict WAT not
0.05/1/0	0.214447	5./51054	1,586.0	141.78	LOPOIOgy_TIST_WAI.pst
	0.1364//	0.040596	044.0	1/2.04	TMC NDT inn
		12 646622	176.0	244.67	
	0.098442	12.646622	1/0.0		
	0.078009	11.50/646	25.0		
1				Input	Output analysis
Table 3: I	File Output S	Stats by Filer	name		
Write	Write	Write Rate	Writes		
Time	MBytes	MBytes/sec		Call	PE=HIDE
0.162883	14.490714	88.963763	5,203.0	2,920.36	Total
0.096137	13.861026	144.179480	3,805.0	3,819.80	tmc_traj_T270.xyz
0.021800	0.064217	2.945740	18.0	3,740.89	<pre>tmc_E_worker_1.out</pre>
0.016016	0.064296	4.014441	18.0	3,745.50	<pre>tmc_E_worker_6.out</pre>
0.013735	0.155310	11.307340	761.0	214.00	tmc_traj_T270.cell
0.004775	0.063504	13.300140	18.0	3,699.39	tmc_E_worker_7.out
0.003025	0.026007	8.596676	505.0	54.00	stdout
0.001983	0.064375	32.470347	19.0	3,552.74	<pre>tmc_E_worker_3.out</pre>
0.001915	0.064375	33.624425	19.0	3,552.74	tmc_E_worker_2.out
0.001905	0.063979	33.588895	18.0	3,727.06	tmc_E_worker_4.out
0.001582	0.063504	40.142573	18.0	3,699.39	tmc_E_worker_5.out
0.000011	0.000122	11.053907	4.0	32.00	_UnknownFile_
========					

Rank reorder suggestions

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Steps to Collecting Performance Data (2)

 At this stage the report gives us useful information and we should get sample hits in time-consuming code sections

• We can see more info

- pat_report a.out+20199-40s.ap2
- You should see this printed to stdout
- Includes
 - job info
 - profile by functions
 - observations and suggestions
 - runtime environment variables
 - hardware performance counter events
- We can also view graphically with Apprentice²
- We can go further on to tracing and loop profiling

Example: Tracing report

- Access perftools, then build and run application
 - module load perftools-base
 - module load perftools-lite-event
 - make clean; make
 - aprun ... a.out
- Comparable to sampling experiment, but now the function are really traced from beginning to end
- Again observations and suggestions are printed
 - E.g. rank reordering
 - And IO observations

Real time in functions Synchronization > cat iob.out Table 1: Profile by Function Group and Function (top 4 functions shown) Time% Time Imb. Calls Group Imb. | Time Time% Function PE=HIDE 100.0% 3.075490 562,739.2 |Total 74.2% | 2.282250 9,855.8 |MPI SYNC 50.8% | 1.562708 | 1.551026 | 99.3% 3,131.2 mpi_bcast_(sync) 12.9% 0.396947 0.396920 100.0% 1.0 |mpi_init_(sync) 91.1% 6,721.6 |mpi allred (sync) 10.5% 0.322147 0.293341 _____ 2.0 USER 19.2% | 0.590584 | 0.661898 54.0% 1.0 |main 5.4% 0.166062 -- | 552.576.7 |MPI 4.1% | 0.126472 | 0.779788 | 87.9% | 541,104.1 |MPI_IPROBE **User functions** Communication

17

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Example: Generate a loop Profile

- Access performance tools software, provide basic tools and environment settings
 - module load perftools-base
- Set environment for tracing experiments with loop profiling
 - module load perftools-loops
- Build instrumented version of the application
 - make clean
 - make
 - You should get an instrumented version program a.out
 - This has been instrumented for sampling (automatic profiling analysis), check with
 - strings a.out | grep 'CrayPat/X' CrayPat/X: Version 6.3.0 Revision 14319 09/02/15 13:51:12

• Run application to get top time consuming routines

- aprun ... a.out (or qsub <pat script>)
- You should get *.rpt and a *.ap2 files
- The report in *.rpt is additionally printed to stdout

Example: Generate a loop Profile

	Subro	outine			Line ni	umber
			I			
Table 1: Inclusive and	l Exclusive	Time in L	.oops (f	rom -hpro	ofile_ge	nerate)
Loop Loop Incl	Time	Loop	Loop	Loop	Loop	<pre>Function=/.LOOP[.]</pre>
Incl Time	(Loop	Hit	Trips	Trips	Trips	PE=HIDE
Time%	Adj.)	I	Avg	Min	Max	
93.0% 19.232051	0.000849	2	26.5	3	50	jacobi .100P.1.1i.236
77.8% 16.092021	0.001350	53	255.0	255	255	jacobi.LOOP.2.li.240
77.8% 16.090671	0.110827	13515	255.0	255	255	jacobi.LOOP.3.li.241
77.3% 15.979844 1 1	5 979844	3446325	<u>511 0</u>	<u>511</u>	<u>511</u>	jacobi.LOOP.4.li.242
14.1% 2.906115	0.001238	53	255.0	255	255	jacobi.LOOP.5.li.263
14.0% 2.904878	0.688611	13515	255.0	255	255	jacobi.LOOP.6.li.264
10.7% 2.216267	2.216267	3446325	511.0	511	511	jacobi.LOOP.7.li.265
4.3% 0.881573	0.000010	1	259.0	259	259	initmt.LOOP.1.li.191
4.3% 0.881563	0.000645	259	259.0	259	259	initmt.LOOP.2.li.192
4.3% 0.880918	0.880918	67081	515.0	515	515	initmt.LOOP.3.li.193
2.7% 0.560499	0.000055	1	257.0	257	257	initmt.LOOP.4.li.210
2.7% 0.560444	0.006603	257	257.0	257	257	initmt.LOOP.5.li.211
2.	3842	66049	513.0	513	513	initmt.LOOP.6.li.212
Nested Loop	os 🛛					

perftools-lite vs. perftools

• There are two ways of using CrayPat

• perftools-lite

- An entry-level approach
- Aimed at users unfamiliar with the full perftools framework
- Provides a report automatically at the end of the job
 - Measures the basic set of performance statistics

perftools

- A more advanced environment
- Provides full control over the performance statistics collected
- Requires a few more steps from the user

• Both generate results as:

- a text report
- a data file (*.ap2) that can be explored using a GUI (Cray Apprentice²)

Steps to Collect Performance Data with perftools

- Access performance tools software
 - module load perftools-base
 - module load perftools
- Build application keeping .o files (CCE: -h keepfiles)
 - make clean
 - make
- Instrument application for automatic profiling analysis
 - pat_build -O apa a.out
 - You should get an instrumented program a.out+pat
 - This has been instrumented for sampling
- Run application to get top time consuming routines
 - aprun ... a.out+pat (or qsub <pat script>)
 - You should get one or more *.xf performance files

Steps to Collecting Performance Data with perftools (2)

- Run pat_report, on the .xf file or the directory
 - pat_report -o <report> <xf file>
 - pat_report -o <report> <xf directory>
 - Generates text report and an .apa instrumentation file
 - We'll discuss pat_report in more detail later
- At this stage the report gives us useful information and we should get sample hits in time-consuming code sections
- We use the .apa file to re-instrument binary for tracing
 - the most important functions have been identified for tracing
- We can inspect and edit the .apa file at this point
 - if we want to tweak the choice of routines to be traced

APA File Example

<pre># You can edit this file, if desired, and use it # to reinstrument the program for tracing like this: # pat_build -0 standard.cray-xt.PE-2.1.56HD.pgi-8.0.amd64.pat- 5.0.0.2- Oapa.512.quad.cores.seal.090405.1154.mpi.pat_rt_exp=default.pat_rt_hwpc=no ne.14999.xf.xf.apa # # These suggested trace options are based on data from: #</pre>	<pre># 31.29% 38517 bytes</pre>
<pre># /home/users/malice/pat/Runs/Runs.seal.pat5001.2009Apr04/./pat.quad/homme/s tandard.cray-xt.PE-2.1.56HD.pgi-8.0.amd64.pat-5.0.0.2- Oapa.512.quad.cores.seal.090405.1154.mpi.pat_rt_exp=default.pat_rt_hwpc=no ne.14999.xf.xf.cdb #</pre>	<pre> # 2.95% 3067 bytes</pre>
<pre># HWPC group to collect by defaultDrtenv=PAT_RT_HWPC=1 # Summary with TLB metrics. #</pre>	<pre># Functions below this point account for less than 10% of samples. # 0.66% 4575 bytes # -T bndry_mod_bndry_exchangev_thsave_time_</pre>
<pre># Libraries to traceg mpi # # User defined functions to trace another by % of complete</pre>	<pre># 0.10% 46797 bytes # -T baroclinic_inst_mod_binst_init_state_ # 0.04% 62214 bytes # -T prim_state_mod_prim_printstate_</pre>
<pre># Oser-defined functions to trace, softed by % of samples. # The way these functions are filtered can be controlled with # pat_report options (values used for this file are shown): # # -s apa_max_count=200 No more than 200 functions are listed. # -s apa_min_size=800 Commented out if text size < 800 bytes.</pre>	<pre></pre>
<pre># -s apa_min_pct=1 Commented out if it had < 1% of samples. -s apa_max_cum_pct=90 Commented out after cumulative 90%. # Local functions are listed for completeness, but cannot be traced. -w # Enable tracing of user-defined functions. # Note: -u should NOT be specified as an additional option.</pre>	<pre># New instrumented program. /.AUTO/cray/css.pe_tools/malice/craypat/build/pat/2009Apr03/2.1.56HD/amd64 /homme/pgi/pat-5.0.0.2/homme/2005Dec08/build.Linux/preqx.cray-xt.PE- 2.1.56HD.pgi-8.0.amd64.pat-5.0.0.2.x # Original program.</pre>

23

Effectively a series of command line arguments to pat_build

Generating Event Traced Profile from APA

- Re-instrument application for further analysis
 - pat_build -0 <apa file>
 - creates new binary: <exe>+apa
- Re-run application
 - aprun … a.out+apa (or qsub <apa script>)
 - This generates a new set of .xf data files
- Generate new text report and visualization file (.ap2)
 - pat_report -o <report> <xf file>
 - pat_report -o <report> <xf directory>
- View report in text and/or with Cray Apprentice2
 - app2 <ap2 file>
 - We'll cover this in more detail later

Steps to Using CrayPat with perftools-lite

Access light version of performance tools software

- > module load perftools-base
- > module load perftools-lite

Build program

> make



a.out (instrumented program)

Run program (no modification to batch script)

aprun a.out
Condensed report to stdout
a.out*.rpt (same as stdout)
a.out*.ap2
MPICH_RANK_XXX files



Steps to Using CrayPat "classic" with perftools



CrayPat (perftools) vs CrayPat (perftools-lite)

- Both use the same process under the hood
- With perftools-lite pat_build runs automatically when the code is linked
 - but keeps the same executable name
- The sample_profile is equivalent to
 - pat_build -0 apa a.out
 - CRAYPAT_LITE = sample_profile (perftools-lite)
- The event_profile is equivalent to
 - pat_build -u -gmpi a.out
 - CRAYPAT_LITE = event_profile (perftools-lite-event)
- It also runs pat_report automatically
 - at the end of the job

Analysing Data with pat_report

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Using pat_report

• pat_report converts raw profiling data into a profile

- Combines .xf data with binary
 - Instrumented binary must still exist when data is converted!
- Produces a text report and an .ap2 file
- .ap2 file can be used for further pat_report calls or display in GUI

• Generates a text report of performance results

- Data laid out in tables
- Many options for sorting, slicing or dicing data in the tables.
 - pat_report -0 *.ap2
 - pat_report -0 help (list of available profiles)
- Volume and type of information depends upon sampling vs tracing.

Advantages of the .ap2 file

- .ap2 file is a self contained compressed performance file
 - Normally it is about 5 times smaller than the .xf file
 - Contains the information needed from the application binary
 - Can be reused
- Independent of the perftools version used to generate it
 - The xf files are very version-dependent
- It is the only input format accepted by Cray Apprentice²
- Once you have the .ap2 file, you can delete:
 - the .xf files
 - the instrumented binary

Files Generated and the Naming Convention

File Suffix	Description
a.out+pat	Program instrumented for data collection
a.out…s.xf	Raw data from sampling experiment available after application execution
a.out…t.xf	Raw data from trace (summarized or full) experiment available after application execution
a.outap2	Processed data, generated by pat_report, contains application symbol information
a.out…s.apa	Automatic profiling analysis template, generated by pat_report (based on pat_build -O apa experiment)
a.out+apa	Program instrumented using .apa file
MPICH_RANK_ORDER.Custom	Rank reorder file generated by pat_report from automatic grid detection an reorder suggestions

```
CrayPat/X: Version 5.2.3.8078 Revision 8078 (xf 8063) 08/25/11 ...
Number of PEs (MPI ranks):
                            16
Numbers of PEs per Node: 16
Numbers of Threads per PE: 1
Number of Cores per Socket: 12
Execution start time: Thu Aug 25 14:16:51 201
System type and speed: x86_64 2000 MHz
Current path to data file:
  /lus/scratch/heidi/ted swim/mpi-openmp/run/swim+pat+27472-34t.ap2
Notes for table 1:
...
```

Sampling Output (Table 1)

N	otes for	table 1	•		
•	••				
Т	able 1:	Profile	by Func	tion	
	Samp %	Samp	Imb. Samp	Imb. Group Samp % Function PE='HIDE'	
	100.0%	775		Total	
	94.2%	730	Ī	USER	
	43.4% 16.1% 8.0% 64.9% 7% 1.7% 1.3% 1.0%	336 125 62 53 28 17 13 11 10 88	8.7285 8.22884 6.1.2.95 1	2.6% mlwxyz_ 4.9% half 9.5% full 3.5% artv_ 3.6% bnd 6.9% currenf_ 8.6% bndsf 13.5% model_ 12.2% cfl 7.0% currenh_ 41.9% bndbo_ 53.4% bndto_	
	5.4%	42	<u> </u>	MPI	
	$ 1.9\% \\ 1.8\% \\ 1.7\%$	15 14 13	4.62 16.53 5.66	23.9% mpi_sendrec 55.0% mpi_bcast 30.7% mpi_barrier	v_ _

pat_report: Flat Profile

Table 1: Profile by Fu	nction Group and	d Function	
Time % Time I	mb. Time Im Time 	b. Calls Gr % F 	oup unction PE='HIDE'
100.0% 104.593634		22649 To	tal
71.0% 74.230520		10473 M	PI
69.7% 72.905208 1.0% 1.050931	0.508369 0.030042	0.7% 125 2.8% 94	mpi_allreduce_ mpi_alltoall_
====================================		73 U	======= SER
16.7% 17.461110 7.7% 8.078474	0.329532 0.114913	1.9% 23 1.4% 48	selfgravity_ ffte4_
2.5% 2.659429		435 M	PI_SYNC
2.1% 2.207467	0.768347 2	6.2% 172	mpi_barrier_(sync)
1.1% 1.188998		11608 H	======= EAP
1.1% 1.166707	0.142473 1	1.1% 5235	free

pat_report: Message Stats by Caller

Таb	le 4: MPI	[Message	Stats by	Caller	
	MPI Msg Bytes	MPI Msg Count	MsgSz <16B Count	4KB<= MsgSz <64KB Count	Function Caller PE[mmm]
15 	138076.0	4099.4	411.6	3687.8	Total
	5138028.0	4093.4	405.6	3687.8	MP1_1SENL
- 3	8080500.0) 2062. <u></u> 	5 93.	8 1968.	8 calc2_ MAIN_
4 4 4	8216000 8208000 6160000).0 3000).0 2000).0 2000).0 100).0).0 50	0.0 200 200 0.0 150	0.0 pe.0 0.0 pe.9 0.0 pe.15
Ιİİ	========				==========
 3	6285250.0) 1656.2 	2 125.0	0 1531. 	2 calc1_ MAIN_
411	8216006	0.0 3000		.0 200	0.0 pe.0
4	6156006	0.0 1500	9.0	150	0.0 pe.3
4	6156000	0.0 1500	9.0 	150	0.0 pe.5

Some important options to pat_report -0

callers **Profile by Function and Callers** callers+hwpc **Profile by Function and Callers** callers+src Profile by Function and Callers, with Line Numbers callers+src+hwpc Profile by Function and Callers, with Line Numbers calltree Function Calltree View heap_hiwater Heap Stats during Main Program **Program HW Performance Counter Data** hwpc Load Balance across PEs load balance program+hwpc load balance sm Load Balance with MPI Sent Message Stats loop_times Loop Stats by Function (from -hprofile generate) loops Loop Stats by Inclusive Time (from -hprofile generate) MPI Message Stats by Caller mpi callers profile Profile by Function Group and Function profile+src+hwpc Profile by Group, Function, and Line samp profile **Profile by Function** samp profile+hwpc **Profile by Function** samp profile+src Profile by Group, Function, and Line

• For a full list see: pat_report -0 help

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Loop Statistics

- Just like adding automatic tracing at the function level, we can add tracing to individual loops.
- Helps identify candidates for parallelization:
 - Loop timings approximate how much work exists within a loop
 - Trip counts can be used to understand parallelism potential
 - useful if considering porting to manycore

• Only available with CCE:

- Requires compiler add additional features into the code.
- Should be done as separate profiling experiment
 - compiler optimizations are restricted with this feature

Loop statistics reported by default in pat_report table

Collecting Loop Statistics

- Load PrgEnv-cray module (default on most systems)
- Load perftools module
- Compile AND link with CCE flag: -h profile_generate

Instrument binary for tracing

- All user functions: pat_build -u my_program
- Or even no user functions: pat_build -w my_program
 - This is sufficient for loop-level profiling of all loops!
- Or use an existing apa file.
- Run the application
- Create report with loop statistics
 - pat_report <xf file> > <report file>

Default Report Table 2

Notes for table 2: Table option: -0 loops

The Function value for each data item is the avg of the PE values. (To specify different aggregations, see: pat help report options s1)

This table shows only lines with Loop Incl Time / Total > 0.009 Profile guided (To set thresholds to zero, specify: -T) optimization

Loop instrumentation can interfere with optimizations, so time feedback for reported here may not reflect time in a fully optimized program. Compiler:

Loop stats can safely be used in the compiler directives: !PGO\$ loop_info est_trips(Avg) min_trips(Min) max_trips(Max) #pragma pgo loop info est trips(Avg) min trips(Min) max trips(Max)

Explanation of Loop Notes (P=1 is highest priority, P=0 is lowest): novec (P=0.5): Loop not vectorized (see compiler messages for reason). sunwind (P=1): Loop could be vectorized and unwound. vector (P=0.1): Already a vector loop.

Default Report Table 2

Table 2:	Loop Stats	s from -hpro	file_gene	erate		
Loop	Loop Incl	Loop Incl	Loop	Loop	Loop	Function=/.LOOP\.
Incl	Time	Time /	Hit	Trips	Notes	PE='HIDE'
Time /		Hit	i	Avg	İ	
Total		Í	Í	ĺ	ĺ	
24.6%	0.057045	0.000570	100	64.1	novec	calc2LOOP.0.li.614
24.0%	0.055725	0.000009	6413	512.0	vector	calc2LOOP.1.li.615
18.9%	0.043875	0.000439	100	64.1	novec	calc1LOOP.0.li.442
18.3%	0.042549	0.000007	6413	512.0	vector	calc1LOOP.1.li.443
17.1%	0.039822	0.000406	98	64.1	novec	calc3LOOP.0.li.787
16.7%	0.038883	0.000006	6284	512.0	vector	calc3LOOP.1.li.788
9.7%	0.022493	0.000230	98	512.0	vector	calc3LOOP.2.li.805
4.2%	0.009837	0.000098	100	512.0	vector	calc2LOOP.2.li.640
=======	=============		========	========	========	==============

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Step 4: Assessing the big picture

• Profile = Where the most of the time is really being spent?

- See also the call-tree view
- Ignore (from the optimization point-of-view) user routines with less than 5% of the execution time
- Why does the scaling end: the major differences in these two profiles?
 - Has the MPI fraction 'blown up' in the larger run?
 - Have the load imbalances increased dramatically?
 - Has something else emerged to the profile?
 - Has the time spent for user routines decreased as it should (i.e. do they scale independently)?

Example with CrayPat



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Step 5: Analyze load imbalance

• What is causing the imbalance?

Computation

 Tasks call for computational kernels (user functions, BLAS routines,...) for varying times and/or the execution time varies depending on the input/caller

Communication

• Large MPI_Sync times

• I/O

 One or more tasks are performing I/O and the others are just waiting for them in order to proceed

Example with CrayPat

22 I			
)21P+hycomBase.ap2 💥		Values	
🦰 љ 🗉			
rview 🗶 🖵 Callgraph 🗶 🔽 Load Balance 🗶			
Calls	Load Balance: mpi_waitall_		
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2			
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8e+04	0 0 41 1.2e+02	4	.3e+0

Step 6: Analyze communication

 What communication pattern is dominating the true time spent for MPI (excluding the sync times)

- Refer to the call-tree view on Apprentice2 and the "MPI Message Stats" tables in the text reports produced by pat_report
- Note that the analysis tools may report load imbalances as "real" communication
 - Put an MPI_Barrier before the suspicious routine load imbalance will aggregate into it in when then analysis is rerun

• How does the message-size profile look like?

• Are there a lot of small messages?

Example with CrayPat report (message stats)

Та	b.	le 4: MP	I Mes	sage S	tat	s by C	all	.er		
		MPI Msg Bytes	MPI Co 	Msg unt 	Ms < Co	gSz 16B unt 	4K Ms <6 Co	B<= gSz 4KB ount	Fun Ca P	ction ller E[mmm]
1	15:	138076.0	409	9.4	41	1.6	368	7.8	Tot	al
İ.	1!	5138028.0	40	93.4	4	05.6	36	87.8	MP	I_ISEN[
 3	- ·	8080500.0	0 2	062.5		93.8	1	.968.8	C	alc2_ MAIN_
 4 4 4		821600 820800 616000	0.0 0.0 0.0	3000. 2000. 2000.	0 0 0	1000. 500.	0	2000 2000 1500).0).0).0	pe.0 pe.9 pe.15
 3		======================================	===== 0 1 	====== 656.2		====== 125.0	1	.531.2	2 C	===== alc1_ MAIN_
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Ì	İ	========	=====	=====	===	=====	:===	=====	===	=====

Step 7: Analyze I/O

- Trace POSIX I/O calls (fwrite, fread, write, read,...)
- How much I/O?
 - Do the I/O operations take a significant amount of time?
- Are some of the load imbalances or communication bottlenecks in fact due to I/O?
 - Synchronous single writer
 - Insert MPI_Barriers to investigate this

Step 8: Find single-core hotspots

- Remember: pay attention only to user routines that consume significant portion of the total time
- View the key hardware counters, for example
 - L1 and L2 cache metrics
 - use of vector (SSE/AVX) instructions
 - Computational intensity (= ratio of floating point ops / memory accesses)
- CrayPat has mechanisms for finding "the" hotspot in a routine (e.g. in case the routine contains several and/or long loops)
 - CrayPat API
 - Possibility to give labels to "PAT regions"
 - Loop statistics (works only with Cray compiler)
 - Compile & link with CCE using -h profile_generate
 - pat_report will generate loop statistics if the flag is being enabled

Example with CrayPat

USER / conj_grad_.LOOPS Time% 59.5% Time 73.010370 secs Flat profile data Imb. Time 3.563452 secs Imb. Time% 4.7% 1.383 /sec 101.0 calls. Calls 183909710385 PERF COUNT HW CACHE L1D:ACCESS PERF COUNT HW CACHE L1D: PREFETCH 7706793512 PERF COUNT HW CACHE L1D:MISS 21336476999 . . . HW counter values 1961227352 SIMD FP 256: PACKED DOUBLE 189983282830 cycles 100.0% Time User time (approx) 73.042 secs **CPU CLK** 3.454GHz 9.3%peak(DP) HW FP Ops / User time 969.844M/sec 70839736685 ops Total DP ops 969.844M/sec 70839736685 ops Computational intensity 0.37 ops/cycle 0.33 ops/ref MFLOPS (aggregate) 124140.04M/sec 1058.97 refs/miss 2.068 avg uses **TLB** utilization Derived D1 cache hit, miss ratios 90.0% hits 10.0% misses metrics D1 cache utilization (misses) 9.98 refs/miss 1.248 avg hits D2 cache hit, miss ratio 82.5% misses 17.5% hits D1+D2 cache hit, miss ratio 91.7% hits 8.3% misses D1+D2 cache utilization 12.10 refs/miss 1.512 avg hits D2 to D1 bandwidth 18350.176MB/sec 1405449334558 bytes Average Time per Call 0.722875 secs

Example with CrayPat

Loop	Loop Incl	Loop Incl	Loop	Loop	Loop	Function=/.LOOP\.
Incl	Time	Time /	Hit	Trips	Notes	PE='HIDE'
ime /		Hit		Avg		
Total						
24.6%	0.057045	0.000570	100	64.1	novec novec	calc2LOOP.0.1i.6
24.0%	0.055725	0.000009	6413	512.0	vector	calc2LOOP.1.li.6
18.9%	0.043875	0.000439	100	64.1	novec	<pre>calc1LOOP.0.li.4</pre>
18.3%	0.042549	0.000007	6413	512.0	vector	<pre>calc1LOOP.1.li.4</pre>
17.1%	0.039822	0.000406	98	64.1	novec	<pre>calc3LOOP.0.li.7</pre>
16.7%	0.038883	0.000006	6284	512.0	vector	calc3LOOP.1.li.7
9.7%	0.022493	0.000230	98	512.0	vector	calc3 .LOOP.2.li.8
4.2%	0.009837	0.000098	100	512.0	vector	calc2 .LOOP.2.li.6
=======	===========	==============	========	========	==========	
						C

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Hardware Counter Selection

HW counter collection enabled

• export PAT_RT_PERFCTR= <group> | <event list>

<pre>\$> man hwpc Table 5. Intel Haswell Event Sets Group Description 0 D1 with instruction counts 1 Summary with cache and TLB metrics 2 D1, D2, and L3 metrics 6 Micro-op queue stalls 7 Back-end stalls 8 Instructions and branches 9 Instruction cache 10 Cache hierarchy 19 Prefetches 23 Summary with cache and TLB metric</pre>	<pre>\$> papi_avail PAPI Preset Events Name Code Avail Deriv Description (Note) PAPI_L1_DCM 0x8000000 Yes No Level 1 data cache misses PAPI_L1_ICM 0x80000001 Yes No Level 1 inst cache misses PAPI_L2_DCM 0x80000002 Yes Yes Level 2 data cache misses PAPI_FP_OPS 0x80000066 Yes Yes Floating point operations</pre>
	• more details:
1	not holp countons
	pat_neip counters
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Further Information

<u>File Edit View History Bookmarks Tools H</u> elp	
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	Contact us Powered by DITAweb
SUBJECTS / SEARCH	CONTENT
Subjects Search	▲ ► IL X. *****
By Subject 🔹	Performance Analysis
Perfomance Tools × S-2376 × ×	The performance analysis process consists of three basic steps.
Software	1. Instrument the program, to specify what kind of data is to be collected under what conditions.
Perfomance Tools	2. Execute the instrumented program, to generate and capture the desired data.
Software Publication Number	3. Analyze the resulting data.
∘ S-2376 2	The Cray Performance Measure and Analysis Tools (CrayPat) suite consists of the following major components:
	 CrayPat-lite, a simplified and easy-to-use version of CrayPat that provides basic performance analysis information automatically, with a minimum of user interaction. For more information about using CrayPat-lite, see CrayPat-lite.
SUBJECT PUBLICATIONS	CrayPat, the full-featured program analysis tool set. CrayPat in turn consists of the following major components.
Showing results 1 to 2 of 2	• pat_build, the utility used to instrument programs
Performance Measurement and Analysis Tools S-2376-63	• the CrayPat run time environment, which collects the specified performance data during program execution
Performance Measurement and Analysis Tools	 pat_report, the first-level data analysis tool, used to produce text reports or export data for more sophisticated analysis
S-2376-631	• Cray Apprentice2, the second-level data analysis tool, used to visualize, manipulate, explore, and compare sets of program performance data in a GUI environment.
	 Reveal, the next-generation integrated performance analysis and code optimization tool, which enables the user to correlate performance data captured during program execution directly to the original source, and identify opportunities for further optimization.
TABLE OF CONTENTS	Cray PAPI components, which are support packages for those who want to access performance counters. For more information, see Monitor Performance Counters.
Performance Measurement and	components, which can also be accessed when the papit module is loaded.
	NOTE: The perftools-base and papi modules are mutually exclusive. One or the other may be loaded, but not both at the same time.
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In-depth Analysis: Using Cray Apprentice2	
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The Golden Rules of profiling:

• Profile your code

- The compiler/runtime will <u>not</u> do all the optimisation for you.
- Profile your code yourself
 - Don't believe what anyone tells you. They're wrong.
- Profile on the hardware you want to run on
 - Don't profile on your laptop if you plan to run on a Cray system

• Profile your code running the full-sized problem

• The profile will almost certainly be qualitatively different for a test case.

• Keep profiling your code as you optimize

- Concentrate your efforts on the thing that slows your code down.
- This will change as you optimise.
- So keep on profiling.

Performance Optimization: Improving Parallel Scalability

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Scalability bottlenecks

- Review the performance measurements (between the two runs)
- Case: user routines scaling but MPI time blowing up
 - Issue: Not enough to compute in a domain
 - Weak scaling could still continue
 - Issue: Expensive (all-to-all) collectives
 - Issue: Communication increasing as a function of tasks

• Case: MPI_Sync times increasing

- Issue: Load imbalance
 - Tasks not having a balanced role in communication?
 - Tasks not having a balanced role in computation?
 - Synchronous (single-writer) I/O or stderr I/O?

Issue: Load imbalances

• Identify the cause

- How to fix I/O related imbalance will be addressed later
- Unfortunately algorithmic, decomposition and data structure revisions are needed to fix load balance issues
 - Dynamic load balancing schemas
 - MPMD style programming
 - There may be still something we can try without code re-design

Consider hybridization (mixing OpenMP with MPI)

- Reduces the number of MPI tasks less pressure for load balance
- May be doable with very little effort
 - Just plug omp parallel do's/for's to the most intensive loops
- However, in many cases large portions of the code has to be hybridized to outperform flat MPI

Issue: Point-to-point communication consuming time

- Message transfer time \propto latency + message size / bandwidth
 - Latency: Startup for message handling
 - Bandwidth: Network BW / number of messages using the same link
- Reduce latency by aggregating multiple small messages if possible
 - Do not pack manually but use MPI's user-defined datatypes
 - Always use the least general datatype constructor possible

• Bandwidth and latency depend on the used protocol

- Eager or rendezvous
 - Latency and bandwidth higher in rendezvous
- Rendezvous messages usually do not allow for overlap of computation and communication (see the extra slides for explanation), even when using non-blocking communication routines
- The platform will select the protocol basing on the message size, these limits can be adjusted STORE ANALYZE

EAGER potentially allows overlapping



Data is pushed into an empty buffer(s) on the remote processor.

Data is copied from the buffer into the real receive destination when the wait or waitall is called.

Involves an extra memcopy, but much greater opportunity for overlap of computation and communication.

Further info

ANALYZE

RENDEZVOUS does not usually overlap



Time

With rendezvous data transfer is often only occurs during the Wait or Waitall statement.

When the message arrives at the destination, the host CPU is busy doing computation, so is unable to do any message matching.

Control only returns to the library when MPI_Waitall occurs and does not return until all data is transferred.

There has been no overlap of computation and communication.

Further info

ANALYZE

Issue: Point-to-point communication consuming time

- One way to improve performance is to send more messages using the eager protocol
 - This can be done by raising the value of the eager threshold, by setting environment variable: export MPICH_GNI_MAX_EAGER_MSG_SIZE=X
 - Values are in bytes, the default is 8192 bytes. Maximum size is 131072 bytes (128KB)
- Try to post MPI_Irecv calls before the MPI_Isend calls to avoid unnecessary buffer copies
- On Cray XE & XC: Asynchronous Progress Engine
 - Progresses also rendezvous messages on the background by launching an extra helper thread to each MPI task
 - Consult 'man mpi' and there the variable MPICH_NEMESIS_ASYNC_PROGRESS

Issue: Point-to-point communication consuming time

- Minimize the data to be communicated by carefully designing the partitioning of data and computation
 - Example: domain decomposition of a 3D grid (n x n x n) with halos to be communicated, cyclic boundaries



1D decomposition ("slabs"): communication \propto n² * w * 2

2D decomposition ("tubes"): communication $\propto n^2 * p^{-1/2} * w * 4$

3D decomposition ("cubes"): communication $\propto n^2 * p^{-2/3} * w * 6$

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w = halo width p = number of MPI tasks

Issue: Expensive collectives

- Reducing MPI tasks by mixing OpenMP is likely to help
- See if every all-to-all collective operation needs to be allto-all rather than one-to-all or all-to-one
 - Often encountered case: convergence checking
- See if you can live with the basic version of a routine instead of a vector version (MPI_Alltoallv etc)
 - May be faster even if some tasks would be receiving data never referenced

• The MPI 3.0 introduces non-blocking collectives (MPI_lalltoall,...)

- Allow for overlapping collectives with other operations, e.g. computation, I/O or other communication
- Are faster (at least on Cray) than the blocking corresponds even without the overlap, and replacement is trivial

Issue: Expensive collectives

 Hand-written RDMA collectives may outperform those of the MPI library

- Fortran coarrays, Unified Parallel C, MPI one-sided communication
- On Cray XE and XC systems, the sc. DMAPP collectives will (usually significantly) improve the performance of the expensive collectives
 - Enabled by the variable: export MPICH_USE_DMAPP_COLL=1
 - Can be used selectively, e.g. export MPICH_USE_DMAPP_COLL=mpi_allreduce
 - Features some restrictions and requires explicit linking with the corresponding library and using the huge pages; consult 'man mpi'

Issue: Performance bottlenecks due to I/O

• Parallelize your I/O !

- MPI I/O, I/O libraries (HDF5, NetCDF), hand-written schmas,...
- Without parallelization, I/O will be a scalability bottleneck in every application

• Try to hide I/O (asynchronous I/O)





- Available on MPI I/O (MPI_File_iwrite/read(_at))
- One can also add dedicated "I/O servers" into code: separate MPI tasks or dedicating one I/O core per node on a hybrid MPI+OpenMP application

Issue: Performance bottlenecks due to I/O

• Tune filesystem (Lustre) parameters

- Lustre stripe counts & sizes, see "man lfs"
- Rule of thumb:
 - # files > # OSTs => Set stripe_count=1 You will reduce the lustre contention and OST file locking this way and gain performance
 - #files==1 => Set stripe_count=#OSTs
 Assuming you have more than 1 I/O client
 - #files<#OSTs => Select stripe_count so that you use all OSTs

• Use I/O buffering for all sequential I/O

- IOBUF is a library that intercepts standard I/O (stdio) and enables asynchronous caching and prefetching of sequential file access
- No need to modify the source code but just
 - Load the module iobuf
 - Rebuild your application

Issue: Performance bottlenecks due to I/O

- When using MPI-I/O and making non-contiguous writes/reads (e.g. multi-dimensional arrays), always define file views with suitable user-defined types and use collective I/O
 - Performance can be 100x compared to individual I/O





Concluding remarks

- Apply the scientific method to performance engineering: make hypotheses and measurements!
- Scaling up is the most important consideration in HPC
- Possible approaches for alleviating typical scalability bottlenecks
 - Find the optimal decomposition & rank placement
 - Overlap computation & communication use non-blocking communication operations for p2p and collective communication both!
 - Make more messages 'eager' and/or employ the Asynchronous Progress Engine (on Cray)
 - Hybridize (=mix MPI+OpenMP) the code to improve load balance and alleviate bottleneck collectives

• Mind your I/O!

- Use parallel I/O
- Tune filesystem parameters